

# CPU Instructions and Procedure Calls

**CS 154: Computer Architecture**

**Lecture #5**

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# Administrative

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- Lab 02 – due today!
- Lab 03 – stay tuned...

# Lecture Outline

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- MIPS instruction formats
- Refresher on some other MIPS instructions and concepts

*Reference material from CS64 – I'll be going over this a little fast...*

# Example

*What does this do?*



```
.data
name: .asciiz "Lisa speaks "
rtn: .asciiz " languages!\n"
age: .word 7
```

```
.text
```

```
main:
```

```
li $v0, 4
la $a0, name # la = load memory address
syscall
```

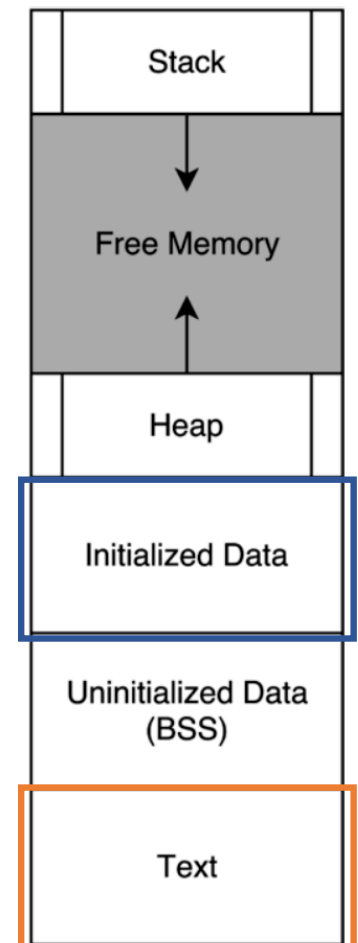
```
la $t2, age
lw $a0, 0($t2)
li $v0, 1
syscall
```

```
li $v0, 4
la $a0, rtn
syscall
```

```
li $v0, 10
syscall
```

*What goes in here?* →

*What goes in here?* →



# .data Declaration Types

## *w/ Examples*

```
var1:    .byte 9          # declare a single byte with value 9
var2:    .half 63        # declare a 16-bit half-word w/ val. 63
var3:    .word 9433      # declare a 32-bit word w/ val. 9433
num1:    .float 3.14     # declare 32-bit floating point number
num2:    .double 6.28    # declare 64-bit floating pointer number
str1:    .ascii "Text"   # declare a string of chars
str3:    .asciiz "Text" # declare a null-terminated string
str2:    .space 5        # reserve 5 bytes of space (useful for arrays)
```

*These are now reserved in memory and we can call them up by loading their memory address into the appropriate registers.*

# Integers in MIPS

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## Unsigned 32-bits

- Range is **0 to  $+2^{32} - 1$**  (or +4,294,967,295)
- Remember positional notation!
  - For when converting to decimal – remember LSB is position **0**
  - Example: What is 0x00881257 in decimal?
  - Answer:  $7 + 2^4 + 2^6 + 2^9 + 2^{12} + 2^{19} + 2^{23} = 8,917,591$

# Integers in MIPS

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## Signed (2s Complement) 32-bits

- Range is  $-2^{31}$  to  $+2^{31} - 1$
- Remember the 2s complement formula!
  - Negate all bits and then add 1
  - Example: What is 0xFFFE775C in decimal?
  - Answer: negative 0x000188A4
    - $= - (4 + 2^5 + 2^7 + 2^{11} + 2^{15} + 2^{16})$
    - $= -10,0516$

# Signed Integers in MIPS

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- Some specific numbers

- 0: 0000 0000 ... 0000
- -1: 1111 1111 ... 1111
- Most-negative: 1000 0000 ... 0000
- Most-positive: 0111 1111 ... 1111

- Representing a number using more bits

- You want to preserve the numeric value
- Example: **+6** in 4-bits (0110) becomes 00000110 in 8-bits
- Example: **-6** in 4-bits (1010) becomes 11111010 in 8-bits
- When does this happen in MIPS?
  - Think of I-type instructions



# MIPS Instructions: Syntax

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**<op> <rd>, <rs>, <rt>**

op : operation

rd : register destination

rs : register source

rt : register target

**<op> <rt>, <rs>, immedi**

op : operation

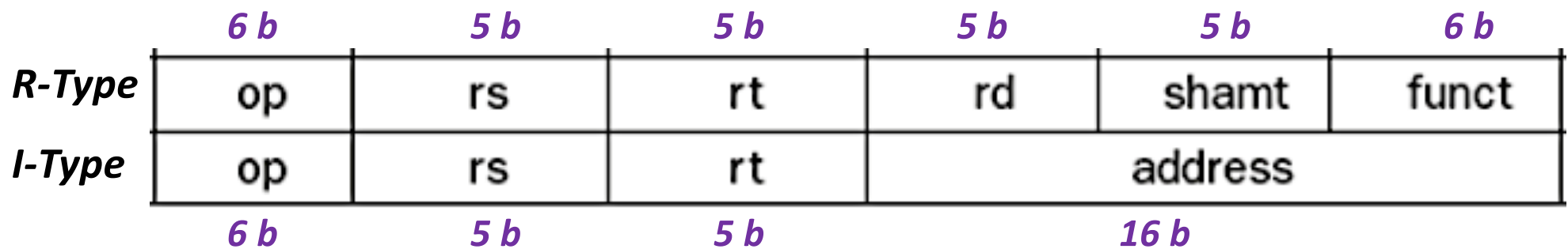
rs : register source

rt : register target

# MIPS Instruction Formats

Recall:

- There are **three** different *instruction formats*: **R**, **I**, **J**
- ALL core instructions are 32 bits long



# Instruction Representation in R-Type

op	rs	rt	rd	shamt	funct
6 b	5 b	5 b	5 b	5 b	6 b
31 - 26	25 - 21	20 - 16	15 - 11	10 - 6	5 - 0

- The combination of the **opcode** and the **funct** code tell the processor what it is supposed to be doing
- Example:

**add \$t0, \$s1, \$s2**

**0x02324020**

op	rs	rt	rd	shamt	funct
0	17	18	8	0	32

op = 0, funct = 32 (0x20)

rs = 17

rt = 18

rd = 8

shamt = 0

means "add"

means "\$s1"

means "\$s2"

means "\$t0"

means this field is unused in this instruction

# Instruction Representation in I-Type

<b>op</b> 6 b 31 – 26	<b>rs</b> 5 b 25 – 21	<b>rt</b> 5 b 20 – 16	<b>address</b> 16 b 15 – 0
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- Example:

**addi \$t0, \$s0, 124**

0x2208007C

<b>op</b> 8	<b>rs</b> 16	<b>rt</b> 8	<b>address/const</b> 124
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op = 8

mean “addi”

rs = 16

means “\$s0”

rt = 8

means “\$t0”

address/const = 124 (0x007C)

is the 16b immediate value

Worth checking out: [https://www.eg.bucknell.edu/~csci320/mips\\_web/](https://www.eg.bucknell.edu/~csci320/mips_web/)

# Pseudoinstructions

- Instructions that are NOT core to the CPU
- They're “macros” of other actual instructions
- Often they are slower (higher CPI) than core instructions

- Examples:

```
li $t0, C
```

Is a macro for:

```
lui $t0, C_hi  
ori $t0, $t, C_lo
```

```
move $t0, $t1
```

Is a macro for:

```
addu $t0, $zero, $t1
```

<https://github.com/MIPT-ILab/mipt-mips/wiki/MIPS-pseudo-instructions> has more examples

# Bitwise Operations

Operation	C/C++	MIPS
Shift left	<<	sll
Shift right	>>	srl, sra
Bitwise AND	&	and, andi
Bitwise OR		or, ori
Bitwise NOT	~	nor*
Bitwise XOR	^	xor

\* Specifically, `nor $t0, $t0, 0` is equivalent to `not(t0)`

# Conditional Operations

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- Branch to a labeled instruction if a condition is true
  - Otherwise, continue sequentially
- **beq rs, rt, L1** often used with **slt**, **slti**
  - if (rs == rt) branch to instruction labeled L1;
- **bne rs, rt, L1** often used with **slt**, **slti**
  - if (rs != rt) branch to instruction labeled L1;
- MIPS also has the pseudoinstructions: **ble**, **blt**, **bge**, **bgt**
  - But pseudoinstructions run slower...
- **j L1**
  - Unconditional jump to instruction labeled L1

# Example

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- C/C++ code: `while (save[i] == k) i += 1;`
- Given: var `i` in `$s3`, `k` in `$s5`, address of `save` in `$s6`
- In MIPS:
  - Loop:

```
sll $t1, $s3, 2
add $t1, $t1, $s6
lw $t0, 0($t1)
bne $t0, $s5, Exit
addi $s3, $s3, 1
j Loop
```
  - Exit: ...



# Procedure Calls (aka Calling Functions)

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- Procedure call: jump and link

**jal FunctionLabel**

- Address of following instruction put in \$ra
- Jumps to target address

- Procedure return: jump register

**jr \$ra**

- Copies \$ra to program counter
- Can also be used for computed jumps
  - e.g., for case/switch statements

# Calling Nested or Recursive Functions

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- What happens when you have a saved return address in \$ra....  
... and then you call ANOTHER function?
- We have to use a standardized way of calling functions
  - The MIPS Calling Convention
- Especially important when different dev. teams are making different functions in a project
  - Also simplifies program testing
- Some registers will be presumed to be “preserved” across a call ;  
Others will not

# The MIPS Calling Convention In Its Essence

- Remember: Preserved vs Unpreserved Regs
  - **Preserved:**                    \$**s0** - \$**s7**, and \$**ra**, and    \$**sp** (by default)
  - **Unpreserved:**    \$**t0** - \$**t9**,    \$**a0** - \$**a3**,    and \$**v0** - \$**v1**
- 
- Values held in **Preserved Regs** immediately before a function call **MUST** be the same immediately after the function returns.
    - Use the **stack memory** to save these
  - Values held in **Unpreserved Regs** must always be assumed to change after a function call is performed.
    - \$**a0** - \$**a3** are for passing arguments into a function
    - \$**v0** - \$**v1** are for passing values from a function

# YOUR TO-DOs for the Week

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- Readings!
  - Chapters 2.10 – 2.13
  
- Stay Tuned for Lab Assignment!
  - Will be announced on Piazza

**</LECTURE>**